

# CS590 Homework 4

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Problem 1:

- (1) Part 1: 4/4. I set up  $A_v$ , decompose  $f|_v = (Af)(v) \cdot \mathbf{1} + g_v$  and apply the variational characterization in the links. Nothing special here.
- (2) Part 2: 1/1. Just induction.
- (3) Part 3: 0.5/1. Minor reasoning mistake handling  $i = k - 1$  but otherwise correct argument.

Problem 2:

- (1) Part 1: 1.5/3. Oversimplified the argument for the second term by using a triangle-ineq type of argument; what actually goes on is a majority-vs-sample comparison. The control over  $Q_3$  is somewhat flawed too by the same token.
- (2) Part 2: 3/3 skipped.

Overall: 10/12.

*Solution to problem 1. Part 1.* We follow the base-case proof of the trickle-down theorem as in lecture 10 and track the smallest eigenvalue. For each  $v \in X(0)$ , let  $A_v$  be the random walk on the 1-skeleton of the link  $X_v$ . Recall that (i) for every  $f : X(0) \rightarrow \mathbb{R}$ ,

$$\langle f, Af \rangle = \mathbb{E}_{v \sim \pi_0} [\langle f|_v, A_v f|_v \rangle_v]$$

and (ii) if we decompose  $f|_v = (Af)(v) \cdot \mathbf{1} + g_v$  with  $g_v \perp \mathbf{1}$  in the link, then

$$\mathbb{E}_{v \sim \pi_0} \|f|_v\|_v^2 = \|f\|^2, \quad \text{and} \quad \langle f|_v, A_v f|_v \rangle_v = \langle g_v, A_v g_v \rangle_v + (Af)(v)^2.$$

Following the theorem assumption, let  $\alpha = \lambda_{\min}(A)$  and pick an eigenfunction  $f$  with  $Af = \alpha f$ . Using above and the variational characterization of the smallest eigenvalue in each link,

$$\alpha \|f\|^2 = \langle f, Af \rangle = \mathbb{E}_v [\langle g_v, A_v g_v \rangle_v + (Af)(v)^2] \geq \lambda_{\min} \mathbb{E}_v \|g_v\|_v^2 + \|Af\|^2.$$

Since  $\|f|_v\|_v^2 = \|g_v\|_v^2 + (Af)(v)^2$  and  $\mathbb{E}_v \|f|_v\|_v^2 = \|f\|^2$ , we have  $\mathbb{E}_v \|g_v\|_v^2 = \|f\|^2 - \|Af\|^2$ . Therefore

$$\alpha \|f\|^2 \geq \lambda_{\min} (\|f\|^2 - \alpha^2 \|f\|^2) + \alpha^2 \|f\|^2,$$

from which the result follows from dividing by  $\|f\|^2 > 0$ ,  $1 - \lambda_{\min} > 0$ , and rearranging.

**Part 2.** We induct on  $m = d - i - 1$  where  $i < d - 2$ . When  $m = 1$ , the 1-skeleton is a graph, so the random walk matrix has  $\lambda_{\min} \geq -1$ . From step  $m$  to  $m + 1$ , we let  $Y$  be any  $(m + 1)$ -dimensional complex. For every  $v \in Y(0)$ ,

the link  $Y_v$  is  $m$ -dimensional, so by IH,  $\lambda_{\min}(Y_v^{\leq 1}) \geq -1/m$ . Applying part 1 to the 2-skeleton  $Y^{\leq 2}$  yields the inductive step:

$$\lambda_{\min}(Y^{\leq 1}) \geq \frac{-1/m}{1 - (-1/m)} = -\frac{1}{m+1}.$$

**Step 3.** Let  $1 \leq k \leq d$ . For any  $i \leq k-2$  and  $\sigma \in X^{\leq k}(i)$ , the 1-skeleton of  $(X^{\leq k})_\sigma$  equals that of  $X_\sigma$ . From lecture 10, every link graph  $X_\sigma$  has  $\lambda_2(X_\sigma) \leq \lambda$  since  $X$  is a  $d$ -dimensional one-sided  $\lambda$ -expander. On the lower side, by part 2 applied to  $X$ ,

$$\lambda_{\min}(X_\sigma) \geq -\frac{1}{d-i-1} > -\frac{1}{d-k+1}.$$

Aggregating both bounds gives the desired claim.

*Solution to problem 2.* **Step 1.** Let  $\epsilon = \mathbb{P}_{(s_1, s_2) \sim D}[f_{s_1} |_{s_1 \cap s_2} \neq f_{s_2} |_{s_1 \cap s_2}]$  and let the lemma's constant be  $C > 0$ , i.e.,

$$\mathbb{E}_t[\mathbb{P}_{s \supset t}[f_s |_{s \setminus t} \neq g_{-t} |_{s \setminus t}]] \leq C\epsilon.$$

Define the three (LHS) quantities to be minimized as  $Q_1(t_1)$ ,  $Q_2(t_1, t_2)$ , and  $Q_3(t_1, t_2)$ , respectively. We show that  $\mathbb{E}[Q_1] \leq C\epsilon$ ,  $\mathbb{E}[Q_2] \leq 2C\epsilon$ , and  $\mathbb{E}[Q_3] \leq 2\epsilon$ . Choosing  $(t_1, t_2)$  to minimize  $Q_1, Q_2, Q_3$  then gives  $Q_i \leq 5C\epsilon = \mathcal{O}(\epsilon)$ .

The first term is already done. For  $Q_2$ , by triangle inequality inside the probability,

$$Q_2(t_1, t_2) \leq \mathbb{P}_{s \supset t_1}[g_{-s \setminus t_1} |_{t_1} \neq f_s |_{t_1}] + \mathbb{P}_{s \supset t_1}[f_s |_{t_1} \neq g_{-t_2} |_{t_1}].$$

The first term can be averaged over  $t_1$ ; using bijection between  $t_1 \subset s$  and  $t = s \setminus t_1$  gives

$$\mathbb{E}_{t_1} \mathbb{P}_{s \supset t_1}[g_{-s \setminus t_1} |_{t_1} \neq f_s |_{t_1}] = \mathbb{E}_t \mathbb{P}_{s \supset t}[g_{-t} |_{s \setminus t} \neq f_s |_{s \setminus t}] \leq C\epsilon$$

whereas for the second term, averaging over disjoint  $t_1, t_2$  and using  $s = t_1 \cup t_2$  gives a similar bound. Hence  $\mathbb{E}[Q_2] \leq 2C\epsilon$ . The third term can be bounded likewise. Finally, because  $\mathbb{E}[Q_1 + Q_2 + Q_3] \leq 5C\epsilon = \mathcal{O}(\epsilon)$ , there exist  $t_1, t_2$  satisfying this property.

**Step 2.** From the hint,

$$\mathbf{1}\{f_s \neq g |_{s}\} \leq 2\mathbb{P}_{t \subset s, |t|=k/2}[f_s |_{t} \neq g |_{t}].$$

Averaging over uniform  $s$ , then swapping the sampling order (first let  $t \sim \binom{[n]}{k/2}$ , then letting  $s \sim$  uniform among the  $\binom{[n]}{k}$  supersets of  $t$ ) gives

$$\mathbb{P}_s[f_s \neq g |_{s}] \leq 2\mathbb{E}_t \mathbb{P}_{s \supset t}[f_s |_{t} \neq g |_{t}].$$

Now we fix  $t$  and write  $X = s \setminus t$ , so  $|X| = k/2$ . For this specific pair  $(s, t)$ , basic set theoretic operations give

$$\mathbf{1}\{f_s |_{t} \neq g |_{t}\} \leq \mathbf{1}\{f_s |_{t} \neq g_{-X} |_{t}\} + \mathbf{1}\{g_{-X} |_{t \setminus t_1} \neq g_{-t_1} |_{t \setminus t_1}\} + \mathbf{1}\{g_{-X} |_{t \cap t_1} \neq g_{-t_2} |_{t \cap t_1}\}.$$

We now take expectations of the three terms citing part 1.

For the first term, by symmetry of joint sampling,

$$\mathbb{E}_t \mathbb{P}_{s \supset t}[f_s |_{t} \neq g_{-X} |_{t}] = \mathbb{E}_t \mathbb{P}_{s \supset t}[f_s |_{s \setminus t} \neq g_{-t} |_{s \setminus t}] \leq \mathcal{O}(\epsilon)$$

by the lemma. For the second term, since disagreement on subset  $t \setminus t_1$  is dominated by disagreement on the whole  $s \setminus t_1$ , by part 1, we have

$$\mathbb{E}_t \mathbb{P}_{s \supset t}[g_{-X} |_{t \setminus t_1} \neq g_{-t_1} |_{t \setminus t_1}] \leq \mathbb{P}_{s \supset t_1}[g_{-(s \setminus t_1)} |_{s \setminus t_1} \neq g_{-t_1} |_{s \setminus t_1}] \leq \mathcal{O}(\epsilon).$$

Finally, the last term can be bounded analogously: disagreement on  $t \cap t_1$  is dominated by that on all of  $t_1$ :

$$\mathbb{E}_t \mathbb{P}_{s \supset t} [g_{-X} |_{t \cap t_1} \neq g_{-t_2} |_{t \cap t_1}] \leq \mathbb{P}_{s \supset t_1} [g_{-(s \setminus t_1)} |_{t_1} \neq g_{-t_2} |_{t_1}] \leq \mathcal{O}(\epsilon),$$

by part 1 again. Combining the three inequalities completes the proof.