

## Homework 4 Solution

Release date: Nov. 1, Due date: Nov. 8.

**Guidelines:** Submit your self-grade in pdf format on Gradescope by **5pm on Nov. 8**. Please compare your solutions with the posted ones and assign yourself points honestly:

- Full credit: essentially correct
- 50% credit: good progress but errors/incomplete
- 25% credit: some relevant ideas only
- No credit: no real attempt

If you take partial credit, briefly note why.

**Q 1** (From one-sided HDXs to two-sided HDXs, 6 pts). In this problem we will prove that given a  $d$ -dimensional one-sided  $\lambda$ -expander  $X$  and an integer  $1 \leq k < d$ , its  $k$ -skeleton  $X^{\leq k}$ , which is obtained by removing from  $X$  all faces of dimension  $> k$  and keeping the same distributions over the other faces, is a  $k$ -dimensional two-sided  $\gamma$ -expander where  $\gamma = \max(\lambda, \frac{1}{d-k+1})$ . We break the proof in two steps.

1. (4 pts) First show a variant of the trickle-down theorem:

**Theorem 0.1.** Let  $X$  be a 2-dimensional simplicial complex with weight  $w_2$ . Suppose that for every  $v \in X(0)$ , the 1-skeleton of the link  $X_v$  has its smallest eigenvalue  $\lambda_{\min}(X_v) \geq \lambda_{\min}$ , then the 1-skeleton of  $X$  has its smallest eigenvalue  $\lambda_{\min}(X^{\leq 1}) \geq \frac{\lambda_{\min}}{1-\lambda_{\min}}$ .

2. (1 pt) Use the theorem above to prove that in a  $d$ -dimensional  $X$ , for any  $i < d - 2$  and  $\sigma \in X(i)$ , the 1-skeleton of the link  $X_\sigma$  has its smallest eigenvalue  $\geq -\frac{1}{d-i-1}$ .

(Note that the smallest eigenvalue of any random walk matrix is at least  $-1$ .)

3. (1 pt) Finish the proof that  $X^{\leq k}$  is a  $k$ -dimensional two-sided  $\max(\lambda, \frac{1}{d-k+1})$ -expander.

### Solution

1. Let  $X$  be a 2-dimensional simplicial complex with vertex weights induced by  $w_2$ . Suppose that for every  $v \in X(0)$ , the 1-skeleton of the link  $X_v$  has smallest eigenvalue  $\lambda_{\min}(X_v) \geq \lambda_{\min}$ . Let  $M$  denote the random-walk matrix on the 1-skeleton of  $X$  with stationary distribution  $\pi$ . For any mean-zero function  $f: X(0) \rightarrow \mathbb{R}$ ,

$$\langle f, Mf \rangle_\pi = \mathbb{E}_{v \sim \pi} [\langle f, M_v f \rangle_{\pi_v}],$$

where  $M_v$  is the walk on  $X_v$ . Using  $\lambda_{\min}(M_v) \geq \lambda_{\min}$ , we have

$$\langle f|_v, M_v f|_v \rangle_{\pi_v} \geq \mathbb{E}_{\pi_v} [f|_v]^2 + \lambda_{\min} \cdot (\|f|_v\|_{\pi_v, 2}^2 - \mathbb{E}_{\pi_v} [f|_v]^2) = \lambda_{\min} \|f|_v\|_{\pi_v, 2}^2 - (1 - \lambda_{\min})(Mf(v))^2.$$

Averaging over  $v$  and rearranging gives

$$\langle f, Mf \rangle_\pi \geq \frac{\lambda_{\min}}{1 - \lambda_{\min}} \langle f, f \rangle_\pi.$$

Hence

$$\lambda_{\min}(X^{\leq 1}) \geq \frac{\lambda_{\min}}{1 - \lambda_{\min}}.$$

2. For a  $d$ -dimensional  $X$  and  $\sigma \in X(i)$ , the link  $X_\sigma$  has dimension  $d - i - 1$ . Set  $a_t =$  smallest eigenvalue for links of dimension  $t$ . From (1),  $a_{t+1} \geq \frac{a_t}{1-a_t}$  and  $a_1 = -1$ , so inductively

$$a_t \geq -\frac{1}{t}.$$

Thus for  $i < d - 2$ ,

$$\lambda_{\min}(X_\sigma^{\leq 1}) \geq -\frac{1}{d-i-1}.$$

3.  $X$  is a  $d$ -dimensional one-sided  $\lambda$ -expander, so for all  $\sigma \in X(i)$ ,  $i < d - 1$ ,

$$\lambda_2(X_\sigma^{\leq 1}) \leq \lambda.$$

From (2),  $\lambda_{\min}(X_\sigma^{\leq 1}) \geq -\frac{1}{d-i-1}$ . For the  $k$ -skeleton  $X^{\leq k}$  we only need  $i \leq k - 1$ , and the worst lower bound occurs at  $i = k - 1$ . Hence every relevant link has

$$|\lambda|_2(X_\sigma^{\leq 1}) \leq \max\left(\lambda, \frac{1}{d-k+1}\right),$$

so the two-sided spectral parameter is

$$\gamma = \max\left(\lambda, \frac{1}{d-k+1}\right).$$

□

**Q 2** (Agreement test over the complete complex, 6 pts). In lecture we used the following theorem without giving a proof

**Theorem 0.2.** Let  $n$  be large and  $k$  moderately sized. For any collection of local functions  $\{f_s : s \rightarrow \Sigma\}_{s \in \binom{[n]}{k}}$ , consider the following direct product test  $D$ :

1. Sample a random subset  $t \in \binom{[n]}{k/2}$ .
2. Then independently sample two subsets  $s_1, s_2$  of size  $k$  conditioned on containing  $t$
3. Accept if and only if the two functions agree on their intersection:

$$f_{s_1}|_{s_1 \cap s_2} \equiv f_{s_2}|_{s_1 \cap s_2}.$$

Then for any collection of local functions  $\{f_s : s \rightarrow \Sigma\}$  there exists a global function  $g : [n] \rightarrow \Sigma$  such that

$$\Pr_{s \sim \text{Uniform}(\binom{[n]}{k})} [f_s \neq g|_s] \leq O(1) \cdot \Pr_{(s_1, s_2) \sim D} [f_{s_1}|_{s_1 \cap s_2} \neq f_{s_2}|_{s_1 \cap s_2}].$$

In this problem we will prove this theorem assuming a key lemma. To statement the lemma, we first define for every subset  $t$  of size  $k/2$  an intermediate function

$$g_{-t} : [n] \setminus t \rightarrow \Sigma, \quad g_{-t}(i) = \text{Majority}\{f_s(i) \mid s \supset t\}.$$

The key lemma states that

**Lemma.**

$$\mathbb{E}_{t \in \binom{[n]}{k/2}} \left[ \Pr_{\substack{s \supset t, \\ s \in \binom{[n]}{k}}} [f_s|_{s \setminus t} \neq g_{-t}|_{s \setminus t}] \right] \leq O(1) \cdot \Pr_{(s_1, s_2) \sim D} [f_{s_1}|_{s_1 \cap s_2} \neq f_{s_2}|_{s_1 \cap s_2}].$$

Use this lemma to finish the proof of the theorem in the following two steps.

1. (3 pts) Prove that there exist two disjoint size- $k/2$  sets  $t_1, t_2$  such that

$$\Pr_{s \supset t_1} [g_{-t_1}|_{s \setminus t_1} \neq f_s|_{s \setminus t_1}] \leq O(1) \cdot \Pr_{(s_1, s_2) \sim D} [f_{s_1}|_{s_1 \cap s_2} \neq f_{s_2}|_{s_1 \cap s_2}],$$

and

$$\Pr_{s \supset t_1} [g_{-t_2}|_{t_1} \neq g_{-s \setminus t_1}|_{t_1}] \leq O(1) \cdot \Pr_{(s_1, s_2) \sim D} [f_{s_1}|_{s_1 \cap s_2} \neq f_{s_2}|_{s_1 \cap s_2}],$$

and

$$\Pr_{t \in \binom{[n] \setminus (t_1 \cup t_2)}{k/2}} [g_{-t_1}|_t \neq g_{-t_2}|_t] \leq O(1) \cdot \Pr_{(s_1, s_2) \sim D} [f_{s_1}|_{s_1 \cap s_2} \neq f_{s_2}|_{s_1 \cap s_2}],$$

2. (3 pts) Define  $g : [n] \rightarrow \Sigma$  to be

$$g(i) = \begin{cases} g_{-t_1}(i) & i \notin t_1 \\ g_{-t_2}(i) & i \in t_1 \end{cases}.$$

Prove that

$$\Pr_s [f_s \neq g|_s] \leq O(1) \cdot \Pr_{(s_1, s_2) \sim D} [f_{s_1}|_{s_1 \cap s_2} \neq f_{s_2}|_{s_1 \cap s_2}].$$

**Hint:** If  $f_s \neq g|_s$ , then  $\Pr_{t \subset s} [f_s|_t \neq g|_t] \geq \frac{1}{2}$ .

**Solution**

Write

$$\Delta := \Pr_{(s_1, s_2) \sim D} [f_{s_1}|_{s_1 \cap s_2} \neq f_{s_2}|_{s_1 \cap s_2}].$$

By the lemma there is an absolute constant  $C > 0$  such that

$$\mathbb{E}_{t \in \binom{[n]}{k/2}} \left[ \Pr_{s \supset t} [f_s|_{s \setminus t} \neq g_{-t}|_{s \setminus t}] \right] \leq C\Delta. \quad (1)$$

**1.Existence of disjoint  $t_1, t_2$  with the three bounds.** Define for each  $t \in \binom{[n]}{k/2}$  the quantity

$$A(t) := \Pr_{s \supset t} [f_s|_{s \setminus t} \neq g_{-t}|_{s \setminus t}].$$

By (1) we have  $\mathbb{E}_t [A(t)] \leq C\Delta$ , hence there exists some  $t_1$  with

$$A(t_1) \leq C\Delta. \quad (i)$$

Next we show that we can pick  $t_2$  disjoint from  $t_1$  so that the other two probabilities are also  $\leq O(\Delta)$ . Fix this  $t_1$  and consider  $t_2$  drawn uniformly from  $\binom{[n] \setminus t_1}{k/2}$ . For any fixed  $s \supset t_1$  and any  $i \in t_1$  define the indicator

$$I_{t_2, s, i} := \mathbf{1}\{g_{-t_2}(i) \neq g_{-(s \setminus t_1)}(i)\}.$$

By the definition of majority, for any two index-sets  $u, v$  and any coordinate  $i$ ,

$$\mathbf{1}\{g_{-u}(i) \neq g_{-v}(i)\} \leq \mathbf{1}\{\text{majority over } s \supset u \text{ differs from } f_{s'}(i) \text{ for some } s' \supset v\}$$

and therefore, for every fixed  $s \supset t_1$  and  $i \in t_1$ ,

$$\mathbb{E}_{t_2} [I_{t_2, s, i}] \leq \Pr_{s' \supset t_2} [f_{s'}(i) \neq g_{-t_2}(i)] + \Pr_{s' \supset (s \setminus t_1)} [f_{s'}(i) \neq g_{-(s \setminus t_1)}(i)].$$

Averaging over  $i \in t_1$ , over  $s \supset t_1$ , and then over  $t_2$  gives

$$\mathbb{E}_{t_2} \left[ \Pr_{s \supset t_1} [g_{-t_2}|_{t_1} \neq g_{-(s \setminus t_1)}|_{t_1}] \right] \leq A(t_2) + A(s \setminus t_1)$$

and averaging once more using (1) (applied to  $t_2$  and to  $s \setminus t_1$ ) yields

$$\mathbb{E}_{t_2} \left[ \Pr_{s \supset t_1} [g_{-t_2}|_{t_1} \neq g_{-(s \setminus t_1)}|_{t_1}] \right] \leq 2C\Delta.$$

Hence there exists some  $t_2$  (disjoint from  $t_1$ ) with

$$\Pr_{s \supset t_1} [g_{-t_2}|_{t_1} \neq g_{-(s \setminus t_1)}|_{t_1}] \leq 2C\Delta. \quad (\text{ii})$$

Finally consider the third quantity

$$B(t_1, t_2) := \Pr_{t \in \binom{[n] \setminus (t_1 \cup t_2)}{k/2}} [g_{-t_1}|_t \neq g_{-t_2}|_t].$$

By the same majority-vs-sample argument as above (compare two majorities by comparing to sample values) and averaging over the choices of sets, one gets

$$\mathbb{E}_{t_2} [B(t_1, t_2)] \leq C'\Delta$$

for some absolute constant  $C'$ , so there exists a choice of  $t_2$  (still disjoint from  $t_1$ ) such that

$$B(t_1, t_2) \leq C'\Delta. \quad (\text{iii})$$

Using Markov's inequality we can show that with some blowup in the constants  $2C, C'$  from (ii),(iii), there exists some  $t_2$  satisfying both inequalities. Thus we complete the proof.

**2. The original problem statement has a bug. So please give yourself full credit for this part.** Below we solve a slightly different version of this problem where we pick  $t_1, t_2$  satisfying the following conditions (it's not hard to prove such  $t_1, t_2$  exist).

$$\Pr_{s \supset t_1} [g_{-t_1}|_{s \setminus t_1} \neq f_s|_{s \setminus t_1}] \leq O(\Delta),$$

$$\Pr_{s \supset t_2} [g_{-t_2}|_{s \setminus t_2} \neq f_s|_{s \setminus t_2}] \leq O(\Delta), \text{ and}$$

$$\Pr_{(t, s, s') \sim \mu} [f_{s'}|_{s' \setminus t_1} \neq g_{-t_1}|_{s' \setminus t_1}],$$

where the distribution  $\mu$  is the equal mixture of two distributions  $\mu_1$  and  $\mu_2$  defined as follows:

$\mu_1$ : sample  $(t, s, s') \sim D$  conditioned on that  $s' \supset (t \cap t_1) \cup t_2$ .

$\mu_2$ : sample  $(t, s, s') \sim D$  conditioned on that  $s' \supset t \cup t_1$ .

We shall prove that  $\Pr_s[f_s \neq g|_s] \leq O(\Delta)$ .

Fix a random  $s \in \binom{[n]}{k}$ . If  $f_s \neq g|_s$  then by the hint at least half of the  $\binom{k}{k/2}$  choices of  $t \subset s$  of size  $k/2$  satisfy  $f_s|_t \neq g|_t$ . Therefore

$$\Pr_s[f_s \neq g|_s] \leq 2 \Pr_{(s,t)} [f_s|_t \neq g|_t],$$

where  $t$  is uniform among  $k/2$ -subsets of  $s$ . To bound the RHS we rewrite it as:

$$\begin{aligned} \Pr_{(s,t)} [f_s|_t \neq g|_t] &\leq \Pr_{(s,t)} [f_s|_{t \cap t_1} \neq g_{-t_2}|_{t \cap t_1}] + \Pr_{(s,t)} [f_s|_{t \setminus t_1} \neq g_{-t_1}|_{t \setminus t_1}] \\ &\leq \Pr_{\substack{(s,t,s') \sim D \\ |s' \supset (t \cap t_1) \cup t_2}} [f_s|_{t \cap t_1} \neq f_{s'}|_{t \cap t_1} \vee f_{s'}|_{t \cap t_1} \neq g_{-t_2}|_{t \cap t_1}] \\ &\quad + \Pr_{\substack{(s,t,s') \sim D \\ |s' \supset t \cup t_1}} [f_s|_{t \setminus t_1} \neq f_{s'}|_{t \setminus t_1} \vee f_{s'}|_{t \setminus t_1} \neq g_{-t_1}|_{t \setminus t_1}] \\ &\leq \Pr_{(s,t,s') \sim \mu_1} [f_s|_{s \cap s'} \neq f_{s'}|_{s \cap s'}] + \Pr_{(s,t,s') \sim \mu_1} [f_{s'}|_{s' \setminus t_2} \neq g_{-t_2}|_{s' \setminus t_2}] \\ &\quad + \Pr_{(s,t,s') \sim \mu_2} [f_s|_{s \cap s'} \neq f_{s'}|_{s \cap s'}] + \Pr_{(s,t,s') \sim \mu_2} [f_{s'}|_{s' \setminus t_1} \neq g_{-t_1}|_{s' \setminus t_1}] \\ &\leq O(\Delta) + O(\Delta) + O(\Delta) + O(\Delta), \end{aligned}$$

where the last inequality follows from the conditions satisfied by  $t_1, t_2$ . Therefore

$$\Pr_s[f_s \neq g|_s] \leq O(\Delta).$$

□