

Properties of the companion matrix:

- The minimal polynomial of  $C_f$  is  $f$ ; so is the char poly.
- $C_f$  is cyclic (see below).
- If  $A$  commutes with  $C_f$  then  $A$  is a polynomial of  $C_f$ . Also, the dimension of the centralizer of  $A$ ,  $\{B \in M_n(F) : AB = BA\}$ , is  $n$ .

We say a matrix is **cyclic** iff there exists a column vector  $v$  so that  $v, Av, A^2v, \dots, A^{n-1}v$  is a basis.

**Proposition**

Let  $A$  be  $n \times n$ . TFAE:

- (1)  $A$  is cyclic.
- (2) min poly of  $A$  has degree  $n$  (and equivalently min poly = char poly).

*Proof.* If  $A$  is cyclic then there exists  $v$  with  $v, Av, \dots, A^{n-1}v$  forming a basis. If  $0 \neq h(x)$  for a polynomial of degree  $< n$ , then  $h(A)v \neq 0$ :

$$\begin{aligned}
 h(x) &= h_0 + h_1x + \dots + h_{n-1}x^{n-1} \\
 h(A) &= h_0I + h_1A + \dots + h_{n-1}A^{n-1} \\
 h(A)v &= h_0v + h_1Av + \dots + h_{n-1}A^{n-1}v,
 \end{aligned}$$

and since the  $A^k v$ 's form a basis, the last quantity is 0 only when  $h(x) = 0$ . Hence the min poly must be of degree  $n$ . Converse omitted.  $\square$

For the last claim, suppose  $BA = AB$ . It suffices to show that it is in the span of  $I, A, \dots, A^{n-1}$ . If  $Bv = w$ , then  $BAv = ABv = Aw$  and  $BA^i v = A^i Bv = A^i w$ . Consider  $\theta : B \rightarrow Bv$  from the centralizer to  $V$ . Furthermore it is linear. By rank nullity  $\dim(C(A)) \leq n$  since  $\ker \theta = 0$ . On the other hand,  $\dim(C(A)) \geq n$  since it at least contains  $I, A, \dots, A^{n-1}$  and no polynomial of degree  $n$  can kill them. Thus  $\dim(C(A)) = n$  and  $I, A, \dots, A^{n-1}$  spans the centralizer.

**Remark.** If  $A$  is cyclic then  $A$  is similar to its companion matrix. To see this, we take the basis  $v, Av, \dots, A^{n-1}v$ . Then the matrix w.r.t this matrix is just

$$\begin{bmatrix}
 0 & 0 & \dots & 0 & -b_0 \\
 1 & 0 & \dots & 0 & -b_1 \\
 0 & 1 & \dots & 0 & -b_2 \\
 \vdots & \vdots & \ddots & \vdots & \vdots \\
 0 & 0 & \dots & 1 & -b_{n-1}
 \end{bmatrix}.$$

**Claim:** the char poly is  $x^n b(n-1)x^{n-1} + \dots + b_0$ . It is enough to show that the min poly  $f(x) = x^n + b_{n-1}x^{n-1} + \dots + b_0$ . Note that

$$f(A)v = b_0v + b_1Av + \dots + b_{n-1}A^{n-1}v + A^n v$$

where  $A^n v$ , by the design of the matrix, kills everything before. Hence  $f(A)v = 0$  and it is the lowest degree polynomial killing  $v$ .

### Corollary

If  $A$  is cyclic then  $A$  is similar to  $A^T$ .

*Proof.*  $A$  and  $A^T$  have the same char poly (and the same min poly). □

## 0.1 Polynomials over Fields

Facts about polynomials over fields:

- There exists a division algorithm for  $f(x)/g(x)$  if  $g(x) \neq 0$ .
- There exists  $q(x), r(x)$  such that  $f(x) = q(x)g(x) + r(x)$  with  $\deg r(x) < \deg g(x)$ .

### Corollary: PID

$\mathbb{F}[x]$  is a **principal ideal domain (PID)**, i.e., if  $I$  is an ideal of  $\mathbb{F}[x]$  ( $I$  is closed under addition and pointwise multiplication), then  $I = h(x)\mathbb{F}[x]$  for some  $h(x)$ .

*Proof.* If  $I = 0$  the claim is true. Otherwise we choose  $h(x) \in I$  nonzero with the lowest possible degree with  $h(x)$  monic. Then if  $0 \neq f(x) \in I$ ,  $\deg(f) \geq \deg(h)$ . By division algorithm  $f(x) = h(x)q(x) + r(x)$ . Since  $r(x) \in I$  it must be 0. □

**GCDs of polynomials:** if  $h(x)$  divides  $f(x)$  and  $g(x)$ , then  $h(x)$  divides everything in the ideal

$$f(x)\mathbb{F}[x] + g(x)\mathbb{F}[x] = \ell(x)\mathbb{F}[x] \quad \text{for some } \ell(x).$$

But then  $h = \ell$  up to a scalar multiple. Therefore, indeed we have

$$\gcd(f, g) = a(x)f(x) + b(x)g(x) \quad \text{for some } a(x), b(x).$$

The Euclidean algorithm works almost identically in the polynomial case.